Educa	ation			
Institute of Computing Technology, Chinese Academy of Sciences, Beijing, China09/2014 - 06/2017Master of Engineering, Computer Science09/2014 - 06/2017				
	87.90/100.00, Advised by Prof. Mingyu Chen er thesis: <u>An Application-driven Flow Scheduling in Data Centers</u> (in Chinese)			
Bache	 hong University of Science and Technology, Wuhan, China elor of Engineering, Computer Science 87.66/100; Ranking: 7/320 	09/2010 - 06/2014		
Public	cations			
1.	Zhuang Wang, Ke Liu, Long Li, Weiyi Chen, Mingyu Chen, Lixin Zhang, " <u>A Nov</u> <u>All Routing in All-optical Hypersquare Torus Network</u> ," in <i>Proc. of ACM Inter</i> <i>Computing Frontiers (CF)</i> , 2016.			
2.	Ke Liu, Zhuang Wang , Jack Y. B. Lee, Mingyu Chen, Lixin Zhang, "Adaptive Rate Control over Mobile Data Networks with Heuristic Rate Compensations," in <i>Proc. of IEEE/ACM International Symposium on Quality of Service (IWQoS)</i> , 2016.			
3.				
4.	Zhuang Wang, Weifa Liang, Meitian Huang, Yu Ma, "Delay-Energy Joint Optimization for Task Offloading in Mobile Edge Computing," in CoRRabs/1804.10416 (2018)			
Profes	ssional Experiences			
-	Abdullah University of Science and Technology	01/2019 - present		
Advis	ed by Prof. Marco Canini Scale machine learning with in-network aggregation			
Brown University		08/2018 - 12/2018		
Advis	 Trace packet transmission across different layers in the kernel Measure NIC-to-NIC latencies, the kernel latencies and the application latencies 			
	ralian National University, Canberra, Australia rch Assistant, Advised by Prof. Weifa Liang	08/2017 - 05/2018		

• Designed task offloading algorithms in Mobile Edge Computing

Awards and Honors

Award	Organization	Date
National Scholarship	Ministry of Education of P. R. China	2011
National Encouragement Scholarship	Ministry of Education of P. R. China	2013
National Scholarship	Ministry of Education of P. R. China	2016
International Postgraduate Research Scholarship (IPRS)	Australian government	2017

Selected Projects

*µ*BPF, a performance diagnosis framework in microservices architecture [GitHub]

 μ BPF leverages eBPF to capture TCP metrics from the kernel for performance diagnosis in microservices architectures. It probes two levels of metrics: flows and packets. The flow-level metrics include sixteen elements, such as flight size, CWnd and sampled RTT. The packet-level metrics are the breakdown of RTTs, including latencies in the TCP layer, IP layer, MAC layer and the network (from NIC to NIC). μ BPF also measures the application layer latency. μ BPF solved several challenges, such as network address translation in microservices architecture, clock synchronization and trace sampling.

Haiyun, a labelled network system [GitHub]

Haiyun enables NICs to distinguish packets with different service-level agreements (SLAs). With the support of DPDK and mTCP, Haiyun could significantly reduce the long tail delay in data centers by keeping packets with different SLAs in different queues from NICs to applications.

My work consists of two components. One is a flow generator, which could generate more than 1 million concurrent TCP connections with just one machine. The other is a third-party tool to measure the performance of Haiyun. To measure the delay resulted from packets processing in servers, we use switches' mirror function to monitor packets and measure their delays on a specific server. To reduce the overheads of operating system, we leverage DPDK to bypass the kernel to optimize the performance of this tool.

Panda, a flow scheduling in data centers [GitHub]

Panda emulates shortest job first (SJF) scheduling to optimize the average flow completion time for throughputintensive applications on the premise that flow information is not known *a priori*. Also, it bounds low latency for delay-sensitive applications by leveraging the distinct flow size distributions of the two kinds of applications. At its heart, Panda derives an optimal threshold to divide packets into two categories: large and small, ensuring that small packets dominate traffic from delay-sensitive applications and large ones dominate traffic from throughputintensive applications. Panda allocates each flow a counter initiated with zero. Large packets increase the counter while small packets decrease it. Then Panda assigns priorities to flows according to their counters.